TRIANGULATING POINT SETS

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DEFINITION

A triangulation of a set $P$ of $n$ points in the plane is a graph having $P$ as set of vertices which is rectilinear, planar, and maximal in the number of edges.
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Corollary. All the faces of such a graph are triangles, except for the unbounded one, which is the exterior of the convex hull of $P$. 
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**COMPLEXITY**

Every triangulation of any set $P$ of $n$ points has:

- $2n - h - 2$ triangles
- $3n - h - 3$ edges

where $h$ is the number of vertices of $ch(P)$. 
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**Proof.** Each triangle has exactly 3 edges. Each internal edge belongs to exactly 2 triangles. Each external edge belongs to exactly 1 triangle. Therefore, $3t = 2(e - h) + h = 2e - h$.

According to Euler’s formula: $n + (t + 1) = v + f = e + 2$.

Combining both equations:
- $e = n + t - 1 \Rightarrow 3e = 3n + 3t - 3 = 3n + 2e - h - 3 \Rightarrow e = 3n - h - 3$
- $3t = 2e - h = 6n - 2h - 6 - h = 6n - 3h - 6 \Rightarrow t = 2n - h - 2$
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DEGENERACIES

As you may have noticed, we are assuming that the set $P$ does not contain three or more points on a line. The assumption is hold along the entire chapter.
DATA STRUCTURE

We want to answer the most usual questions for any decomposition of the plane:
- For any given triangle, report its edges/vertices.
- For any given vertex, report the sorted list of edges/triangles incident to it.
- For any given edge, report its endpoints and its adjacent triangles.
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Table of vertices

<table>
<thead>
<tr>
<th>v</th>
<th>x</th>
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DCEL

| e | v_B | v_E | f_L | f_R | e_P | e_N |

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TRIANGULATING POINT SETS

ALGORITHMS
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1. Incremental algorithms
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1. Incremental algorithms
   1.1. Without sorting
1. Incremental algorithms

1.1. Without sorting

For each $i$, detect whether $p_i$ lies in the interior or the exterior of $ch(p_1, \ldots, p_{i-1})$. If it is external, compute the supporting lines from $p_i$ to $ch(p_1, \ldots, p_{i-1})$ and add all the intermediate diagonals to the triangulation. If it is internal, detect the triangle $T$ containing $p_i$ and partition $T$ into 3 triangles.
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   1.1. Without sorting
   1.2. With sorting

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ALGORITHMS

1. Incremental algorithms
   
   1.1. Without sorting

   Running time: \( O(n^2) \)

   1.2. With sorting

   Start by sorting the points in lexicographical order in \( O(n \log n) \) time. The information of the sorted order of the points allows to add the \( i \) diagonals in \( O(i) \) time, so that the amortized cost of the insertion of all diagonals is done in \( O(n) \) time.
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1. Incremental algorithms
   1.1. Without sorting
   1.2. With sorting
   1.3. With hierarchical structure

Running time:

- Without sorting: \( O(n^2) \)
- With sorting: \( O(n \log n) \)
ALGORITHMS

1. Incremental algorithms

1.1. Without sorting  \textbf{Running time: } $O(n^2)$
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Using an auxiliary enclosing triangle and a hierarchy of triangles: each time a new point is added, a triangle gets subdivided into three children.
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1.5. With auxiliary point(s)

A fixed point \( p \) is used as a reference, and \( P \cup \{p\} \) is enclosed in an auxiliary triangle. When inserting each point \( p_i \):
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   1.3. With hierarchical structure
       
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![Diagram of point triangulation](image-url)
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2. Graham’s algorithm
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- **1.3. With hierarchical structure**: $O(n^2)$ worst case, $O(n \log n)$ if balanced
- **1.4. Randomized**: $O(n \log n)$ expected
- **1.5. With auxiliary point(s)**: $O(n^2)$ worst case, $O(n^{3/2})$ expected
ALGORITHMS

1. Incremental algorithms
   1.1. Without sorting
   1.2. With sorting
   1.3. With hierarchical structure
   1.4. Randomized
   1.5. With auxiliary point(s)

2. Graham’s algorithm

Running time: $O(n^2)$

Running time: $O(n \log n)$

Running time: $O(n^2)$ worst case, $O(n \log n)$ if balanced

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ALGORITHMS

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   1.1. Without sorting
   1.2. With sorting
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2. Graham’s algorithm

3. Divide and conquer

Running time:
- $O(n^2)$
- $O(n \log n)$
- $O(n^2)$ worst case, $O(n \log n)$ if balanced
- $O(n \log n)$ expected
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   Running time: $O(n \log n)$

3. Divide and conquer

   Initialization
   Sort the points by abscissa

   Advance
   - Partition: divide the points into roughly two vertically separated halves
   - Recursion: recursively triangulate each half
   - Fusion: compute the external common tangents and triangulate the intermediate space
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   - 1.1. Without sorting
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     - **Running time:** $O(n \log n)$
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   - 1.5. With auxiliary point(s)
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2. Graham’s algorithm
   - **Running time:** $O(n \log n)$

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   **Initialization**
   - Sort the points by abscissa

   **Advance**
   - Partition: divide the points into roughly two vertically separated halves
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TRIANGULATING POINT SETS

ALGORITHMS

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       Running time: $O(n \log n)$
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2. Graham’s algorithm
   Running time: $O(n \log n)$

3. Divide and conquer
   Running time: $O(n \log n)$

LOWER BOUND

This problem has an $\Omega(n \log n)$ lower bound, since the convex hull of the set of points can be trivially obtained in $O(n)$ time from the triangulation.
Quality of a triangulation
Quality of a triangulation
Quality of a triangulation

Incremental, without sorting
Quality of a triangulation

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Graham’s
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Divide and conquer
Quality of a triangulation

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Incremental, unsorted
Incremental, sorted
Graham’s scan
Divide and conquer
Quality of a triangulation

- Incremental, unsorted
- Incremental, sorted
- Graham's scan
- Divide and conquer
- Delaunay triangulation
Quality of a triangulation

Delaunay
Quality of a triangulation

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Delaunay Triangulation
DELAUNAY TRIANGULATION

A TOOL FOR INTERPOLATION
DELAUNAY TRIANGULATION

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DELAUNAY TRIANGULATION

DEFINITION AND PROPERTIES
DELAUNAY TRIANGULATION

DEFINITION AND PROPERTIES

Definition

Given a set $P$ with $n$ points in the plane...
DEFINITION AND PROPERTIES

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Given a set $P$ of $n$ points in the plane, the Delaunay triangulation of $P$, $\text{Del}(P)$, is the rectilinear dual graph of the Voronoi diagram $\text{Vor}(P)$. 
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Characterization

- Two points $p_i, p_j \in P$ form a Delaunay edge if and only if there exists a circle through $p_i$ and $p_j$ which does not contain any point of $P$ in its interior.
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- Three points $p_i, p_j, p_k$ form a Delaunay triangle (in general, are vertices of a face) if and only if the circle through them does not contain any point of $P$ in its interior.
DEFINITION AND PROPERTIES

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Given a set \( P \) of \( n \) points in the plane, the Delaunay triangulation of \( P \), \( Del(P) \), is the rectilinear dual graph of the Voronoi diagram \( Vor(P) \).

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- Two points \( p_i, p_j \in P \) form a Delaunay edge if and only if there exists a circle through \( p_i \) and \( p_j \) which does not contain any point of \( P \) in its interior.

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DELAUNAY TRIANGULATION

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Property
$\text{Del}(P)$ is a plane graph
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If $\overline{pq}$ is a Delaunay edge, there exists an empty circle through $p$ and $q$. If a segment $\overline{rs}$ intersects $\overline{pq}$, then every circle through $r$ and $s$ contains at least one of $p$ or $q$. 
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Property
$\text{Del}(P)$ is a triangulation of $P$, except when $P$ has three or more concyclic points. In this case, it is a pre-triangulation which can be trivally completed (although this can be done in several different ways).
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GLOBAL CHARACTERIZATION

Theorem

\[ T(P) = \text{Del}(P) \text{ iff the circumcircles of the triangles of } T(P) \text{ are empty of points of } P. \]
Theorem

\[ T(P) = Del(P) \text{ iff the circumcircles of the triangles of } T(P) \text{ are empty of points of } P. \]

Let \( p_i \in P \). Let \( p_1, \ldots, p_k \) be the vertices of the triangles of \( T(P) \) incident to \( p_i \), sorted in counterclockwise order, \( C_1, \ldots, C_k \) be their circumcircles, and \( q_1, \ldots, q_k \) their centers (\( q_j \) denotes the center of \( C_j \), the circumcircle of \( p_i, p_j, p_{j+1} \)). We will prove that the polygon \( Q = \{q_1, \ldots, q_k\} \) coincides with \( Vor(p_i) \).
**Theorem**

\[ T(P) = \text{Del}(P) \] iff the circumcircles of the triangles of \( T(P) \) are empty of points of \( P \).

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\[
\overline{q_{j-1}q_j} \perp \overline{p_ip_j} \quad \Rightarrow \quad Q = \bigcap_{j=1}^{k} H_{ij}
\]
Theorem

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\[
\overline{q_{j-1}q_j} \perp \overline{p_ip_j} \implies Q = \bigcap_{j=1}^{k} H_{ij}
\]

If \( h \neq 1, \ldots, k \) then \( q_j \in b(p_i, r_h) \) and, therefore,

\[
\bigcap_{j=1}^{k} H_{ij} \subset H(p_i, r_h) \subset H_{ih}
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Theorem

$T(P) = Del(P)$ iff the circumcircles of the triangles of $T(P)$ are empty of points of $P$.

Let $p_i \in P$. Let $p_1, \ldots, p_k$ be the vertices of the triangles of $T(P)$ incident to $p_i$, sorted in counterclockwise order, $C_1, \ldots, C_k$ be their circumcircles, and $q_1, \ldots, q_k$ their centers ($q_j$ denotes the center of $C_j$, the circumcircle of $p_i, p_j, p_{j+1}$). We will prove that the polygon $Q = \{q_1, \ldots, q_k\}$ coincides with $Vor(p_i)$.

If $h \neq 1, \ldots, k$ then $q_j \in b(p_i, r_h)$ and, therefore,

$$\bigcap_{j=1}^{k} H_{ij} \subset H(p_i, r_h) \subset H_{ih}$$

Hence,

$$Q = \bigcap_{j=1}^{k} H_{ij} = \bigcap_{j \neq i} H_{ij} = Vor(p_i)$$
Theorem

Let $P = \{p_1, \ldots, p_n\}$ with $p_i = (a_i, b_i, 0)$. Let $p_i^* = (a_i, b_i, a_i^2 + b_i^2)$ be the vertical projection of each point $p_i$ onto the paraboloid $z = x^2 + y^2$. Then $\text{Del}(P)$ is the orthogonal projection onto the plane $z = 0$ of the lower convex hull of $P^*$. 
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DELAUNAY TRIANGULATION

DELAUNAY TRIANGULATION AND 3D CONVEX HULL

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Let $P = \{p_1, \ldots, p_n\}$ with $p_i = (a_i, b_i, 0)$. Let $p_i^* = (a_i, b_i, a_i^2 + b_i^2)$ be the vertical projection of each point $p_i$ onto the paraboloid $z = x^2 + y^2$. Then $Del(P)$ is the orthogonal projection onto the plane $z = 0$ of the lower convex hull of $P^*$. 
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Let $P = \{p_1, \ldots, p_n\}$ with $p_i = (a_i, b_i, 0)$. Let $p^*_i = (a_i, b_i, a_i^2 + b_i^2)$ be the vertical projection of each point $p_i$ onto the paraboloid $z = x^2 + y^2$. Then $\text{Del}(P)$ is the orthogonal projection onto the plane $z = 0$ of the lower convex hull of $P^*$. 
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$p_i^*, p_j^*, p_k^*$ form a (triangular) face of the lower convex hull of $P^*$
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Let $P = \{p_1, \ldots, p_n\}$ with $p_i = (a_i, b_i, 0)$. Let $p_i^* = (a_i, b_i, a_i^2 + b_i^2)$ be the vertical projection of each point $p_i$ onto the paraboloid $z = x^2 + y^2$. Then $Del(P)$ is the orthogonal projection onto the plane $z = 0$ of the lower convex hull of $P^*$.

$p_i^*, p_j^*, p_k^*$ form a (triangular) face of the lower convex hull of $P^*$

The plane through $p_i^*, p_j^*, p_k^*$ leaves all the remaining points of $P^*$ above it.
Theorem

Let \( P = \{p_1, \ldots, p_n\} \) with \( p_i = (a_i, b_i, 0) \). Let \( p_i^* = (a_i, b_i, a_i^2 + b_i^2) \) be the vertical projection of each point \( p_i \) onto the paraboloid \( z = x^2 + y^2 \). Then \( \text{Del}(P) \) is the orthogonal projection onto the plane \( z = 0 \) of the lower convex hull of \( P^* \).

\( p_i^*, p_j^*, p_k^* \) form a (triangular) face of the lower convex hull of \( P^* \)

\[ \uparrow \quad \uparrow \]

The plane through \( p_i^*, p_j^*, p_k^* \) leaves all the remaining points of \( P^* \) above it

\[ \uparrow \quad \uparrow \]

The circle through \( p_i, p_j, p_k \) leaves all the remaining points of \( P \) in its exterior
DELAUNAY TRIANGULATION

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The circle through $p_i, p_j, p_k$ leaves all the remaining points of $P$ in its exterior

$p_i, p_j, p_k$ form a triangle of $\text{Del}(P)$
DELAUNAY TRIANGULATION

LOCAL CHARACTERIZATION
A triangulation $T(P)$ is locally Delaunay if each pair of triangles $p_i p_j p_k$ and $p_i p_j p_l$ sharing an edge $p_i p_j$ satisfies $p_l \notin C_{ijk}$ and $p_k \notin C_{ijl}$. 
DELAUNAY TRIANGULATION

LOCAL CHARACTERIZATION

Definition

A triangulation $T(P)$ is **locally Delaunay** if each pair of triangles $p_ip_jp_k$ and $p_ip_jp_l$ sharing an edge $p_ip_j$ satisfies $p_l \notin C_{ijk}$ and $p_k \notin C_{ijl}$.

The edge $p_ip_j$ is locally Delaunay

The edge $p_ip_j$ is not locally Delaunay

The edge $p_ip_j$ is locally Delaunay

In fact, the quadrilateral $p_ip_lp_jp_k$ is not convex.

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Theorem

A triangulation $T(P)$ is a Delaunay triangulation if and only if it is locally Delaunay.
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Let $\overline{p_ip_j}$ be the edge of $T_{ijk}$ separating $p_l$ from $T_{ijk}$. Among all 4-tuples in this situation, let $ijkl$ maximize the angle $p_ip_lp_j$.

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\[\text{Computational Geometry, Facultat d'Informàtica de Barcelona, UPC}\]
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DELAUNAY TRIANGULATION

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As $T(P)$ is locally Delaunay, $m \neq l$.

Then $p_l \in C_{ijm}$.

Hence, one of the angles $p_ip_lp_m$ or $p_jp_lp_m$ would be greater than $p_ip_lp_j$. 
DELAUNAY FLIPS

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**Lemma 1.** Let $C$ be a circle, $\overline{ab}$ a chord of $C$, and $p$, $q$, $r$ and $s$ four points lying to the same side of the line $\overline{ab}$. If $r$ is internal to $C$, $p$ and $q$ lie in $C$, and $s$ is external to $C$, then the following relations hold between the angles formed at $p$, $q$, $r$ and $s$ by the chord $\overline{ab}$: $\hat{s} < \hat{p} = \hat{q} < \hat{r}$. 

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DELAUNAY TRiangulation

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First case:

\[
\begin{align*}
2\delta + 2\gamma + 2\beta &= \pi \\
2\alpha + 2\beta &= \pi
\end{align*}
\]  
\[\Rightarrow 2\alpha = 2\gamma + 2\delta \Rightarrow \alpha = \gamma + \delta \Rightarrow \hat{p} = \hat{q} = \alpha\]
DELAUNAY TRIANGULATION

DELAUNAY FLIPS

We intend to prove that \( \text{Del}(P) \) can be obtained from any triangulation of \( P \) by Delaunay flips, which consist in deleting the diagonal of a convex quadrilateral if it is not locally Delaunay, and replacing it by the other diagonal of the quadrilateral.

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First case:

\[
2\delta + 2\gamma + 2\beta = \pi \quad \text{and} \quad 2\alpha + 2\beta = \pi \quad \Rightarrow \quad 2\alpha = 2\gamma + 2\delta \quad \Rightarrow \quad \alpha = \gamma + \delta \quad \Rightarrow \quad \hat{p} = \hat{q} = \alpha
\]

Second case:

\[
2\alpha + \epsilon + 2\delta = \pi \quad \text{and} \quad 2\gamma + \epsilon = \pi \quad \Rightarrow \quad 2\alpha + 2\delta - 2\gamma = 0 \quad \Rightarrow \quad \alpha = \gamma - \delta \quad \Rightarrow \quad \hat{p} = \hat{q} = \alpha
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We intend to prove that Del(\(P\)) can be obtained from any triangulation of \(P\) by Delaunay flips, which consist in deleting the diagonal of a convex quadrilateral if it is not locally Delaunay, and replacing it by the other diagonal of the quadrilateral.

**Lemma 1.** Let \(C\) be a circle, \(ab\) a chord of \(C\), and \(p, q, r\) and \(s\) four points lying to the same side of the line \(ab\). If \(r\) is internal to \(C\), \(p\) and \(q\) lie in \(C\), and \(s\) is external to \(C\), then the following relations hold between the angles formed at \(p, q, r\) and \(s\) by the chord \(ab\): \(\hat{s} < \hat{p} = \hat{q} < \hat{r}\).

Let us prove that \(\hat{p} = \hat{q}\):

The remaining relations follow immediately:
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We intend to prove that $\text{Del}(P)$ can be obtained from any triangulation of $P$ by Delaunay flips, which consist in deleting the diagonal of a convex quadrilateral if it is not locally Delaunay, and replacing it by the other diagonal of the quadrilateral.

Lemma 2. When the chord $\overline{ab}$ is a diameter of $C$, the angle $\hat{p}$ for any $p \in C$ is $\pi/2$. 
We intend to prove that $\text{Del}(P)$ can be obtained from any triangulation of $P$ by Delaunay flips, which consist in deleting the diagonal of a convex quadrilateral if it is not locally Delaunay, and replacing it by the other diagonal of the quadrilateral.

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Since in this case $2\alpha = \pi$. 

![Diagram of a circle with chord $\overline{ab}$ and Delaunay triangles]

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Lemma 3. Given any chord $\overline{ab}$ in a circle $C$, if one of the arcs corresponds to $\alpha$, then the other one corresponds to $\pi - \alpha$. 
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\[
\begin{align*}
\alpha_1 + \beta + \gamma &= \frac{\pi}{2} \\
\alpha_2 + \beta + \delta &= \frac{\pi}{2}
\end{align*}
\Rightarrow \begin{align*}
\alpha + 2\beta + \gamma + \delta &= \pi \\
x + \gamma + \delta &= \pi \\
2\alpha + 2\beta &= \pi
\end{align*}
\Rightarrow x = \pi - \alpha
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**Lemma 4.** Let $\overline{pq}$ be the common edge of the triangles $pqa$ and $pqb$, forming a convex quadrilateral. Then:

$$a \in \text{ext}(C_{pqb}) \iff b \in \text{ext}(C_{pqa})$$
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**Lemma 5.** Consider a convex quadrilateral with diagonals $\overline{ab}$ and $\overline{pq}$. Then:

\[
\overline{ab} \text{ is not locally Delaunay } \iff \overline{pq} \text{ is locally Delaunay}
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$\overline{ab}$ is not locally Delaunay $\iff$ $\overline{pq}$ is locally Delaunay

$\overline{ab}$ is not locally Delaunay

$\iff q \in \text{int}(\mathcal{C}_{abp})$
$\iff \hat{aqp} > \hat{abp}$
$\iff b \in \text{ext}(\mathcal{C}_{apq})$
$\iff \overline{pq}$ is locally Delaunay
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Lemma 6. Let $P$ be a set of points $p_i = (x_i, y_i, 0)$ in the plane, and let $P^*$ be the set of their vertical projections $p^* = (x_i, y_i, x_i^2 + y_i^2)$ onto the unit paraboloid. Producing a Delaunay flip in a triangulation of $P$ corresponds to “sticking” a tetrahedron from below to the corresponding polyhedrization of $P^*$. 
We intend to prove that $\text{Del}(P)$ can be obtained from any triangulation of $P$ by Delaunay flips, which consist in deleting the diagonal of a convex quadrilateral if it is not locally Delaunay, and replacing it by the other diagonal of the quadrilateral.

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Once flipped, the quadrilateral is locally Delaunay: the fourth point lies in the exterior of the circumcircle of the triangle.

In the paraboloid, this means that the fourth point lies above the triangular face of the polyhedrization.
DELAUNAY FLIPS

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**Corollary.** Given any triangulation of $P$, performing locally Delaunay flips is a procedure converging to $\text{Del}(P)$. 
1. Compute the Voronoi diagram by any of the known methods and dualize it.
ALGORITHMS

1. Compute the Voronoi diagram by any of the known methods and dualize it.

2. Project the points onto the paraboloid, compute the 3D convex hull by any of the known methods, and project it back onto the plane.
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3. Compute a triangulation, by any of the known methods, and apply Delaunay flips.
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3. Compute a triangulation, by any of the known methods, and apply Delaunay flips.

4. Incremental algorithm
   - Compute an enclosing triangle for \( \{p_1, \ldots, p_n\} \)
   - Compute \( \text{Del}(p_1, \ldots, p_{i+1}) \) from \( \text{Del}(p_1, \ldots, p_i) \)
DELAUNAY TRIANGULATION

INCREMENTAL ALGORITHM
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DELAUNAY TRIANGULATION

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Let \( D_i = Del(p_1, \ldots, p_i) \) and \( p = p_{i+1} \).

**Observation 1.** If \( qrs \) is the triangle of \( D_i \) containing \( p \), then \( pq, pr \) and \( ps \) are edges of \( D_{i+1} \).
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Observation 1. If $qrs$ is the triangle of $D_i$ containing $p$, then $pq$, $pr$ and $ps$ are edges of $D_{i+1}$.

As $C_{qrs}$ is empty, there exist empty circles $C_{pq}$, such as the circle through $p$ and $q$ tangent to $C_{qrs}$ in $q$. Similarly for $r$ and $s$. 
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Observation 2. Let \( pqr \) be a triangle incident to \( p \). The edge \( \overline{qr} \) may not be a Delaunay edge.
Let $D_i = Del(p_1, \ldots, p_i)$ and $p = p_{i+1}$.

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Since $p$ may lie in the interior of $C_{qrt}$. 
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Obvious, because the property is local: it affects only quadrilaterals formed by two triangles sharing an edge.
INCREMEITAL ALGORITHM

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Algorithm

Each time a new point is added to the triangulation, and before adding the next point, the following routine is executed:

Flips

While there are still triangles incident to $p$ non locally Delaunay, flip them.
DELAUNAY TRIANGULATION

INCREMENTAL ALGORITHM

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DELTAUNAY TRIANGULATION

INCREMENTAL ALGORITHM

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Each time a new point is added to the triangulation, and before adding the next point, the following routine is executed:

**Flips**

While there are still triangles incident to $p$ non locally Delaunay, flip them.
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The added running time of performing the flips when adding $p_i$ is

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As the average order is smaller than 6, the expected added running time is not $O(n^2)$ but simply $O(n)$. 
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Let us be more precise:

If $\mathcal{T} = \{T_1, \ldots, T_t\}$ is a triangulation of $P$, the “fineness” of $\mathcal{T}$ is the increasingly sorted list of the angles of all the triangles $T_i$ of $\mathcal{T}$: $F(\mathcal{T}) = (\alpha_1, \ldots, \alpha_{3t})$. 
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Since every triangulation of $P$ has $t = 2n - h - 2$ triangles, these $3t$-tuples can be compared and lexicographically sorted.
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The Delaunay triangulation maximizes the “fineness”:

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The proof of this statement requires a last lemma.
Lemma 7. Let $a$, $b$, $c$ and $d$ be four points forming a convex quadrilateral, in counterclockwise order. Let $\mathcal{T}$ and $\mathcal{T}'$ be the two possible triangulations of the quadrilateral: $\mathcal{T}$ uses the diagonal $ac$ and $\mathcal{T}'$ uses $bd$. Let $\epsilon$ and $\epsilon'$ respectively be the minimum angles of $\mathcal{T}$ and $\mathcal{T}'$. Then:

\[
\begin{align*}
\epsilon > \epsilon' & \iff d \in \text{ext}(C_{abc}) \\
\epsilon = \epsilon' & \iff d \in \partial(C_{abc}) \\
\epsilon < \epsilon' & \iff d \in \text{int}(C_{abc})
\end{align*}
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Due to the symmetry of the problem, we only need to prove that $\epsilon > \epsilon' \iff d \in \text{ext}(C_{abc})$. 
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\[ \alpha_1 \quad \beta_1 \quad \alpha_2 \quad \beta_2 \]

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If $P$ contains four or more concyclic points, $\text{Del}(P)$ contains a polygon inscribed in a circle which can be triangulated in several ways. Nevertheless, Lemma 1 (on the geometrical locus of all the points from which a segment is seen under a given angle) guarantees that every triangulation of a polygon inscribed in a circle has the same fineness, since each edge of the polygon belongs to a triangle, and every possible triangle gives rise to the same angle.
The Delaunay triangulation is used to interpolate terrains, because it also minimizes the roughness of the terrain, in other words, the integral of the square of the $L_2$-norm of the terrain’s gradient.

It is important to notice that this property is independent from the data, in other words, it is independent from the values of the $z$-coordinates of the input points.
SOME ADDRESSES TO PLAY WITH DELAUNAY TRIANGULATIONS

http://www.cs.cornell.edu/Info/People/chew/Delaunay.html

http://www.pi6.fernuni-hagen.de/GeomLab/VoroGlide/

http://www.dma.fi.upm.es/docencia/segundociclo/geomcomp/voronoi.html

http://www.cs.unc.edu/~snoeyink/terrain/Demo.html

AND TWO BOOKS WITH MUCH MORE INFORMATION

A. Okabe, B. Boots, K. Sugihara, S. N. Chiu
Spatial Tessellations

F. Aurenhammer, R. Klein, D.-T. Lee
Voronoi Diagrams and Delaunay Triangulations